## Course Overview

1. Biological Background
2. Pairwise sequence alignment algorithms
3. Probabilistic alignments: Hidden Markov models
4. Multiple sequence alignments
5. Phylogeny: Algorithms for reconstructing pedigrees
6. Neural nets \& deep learning for sequence analysis
7. Recent advances \& applications

## Short historical Introduction

- Genetics as a natural science started in 1866: Gregor Mendel performed experiments that pointed to the existence of biological elements called genes.
- Deoxy-ribonucleic acid (DNA) isolated by Friedrich Miescher in 1869.
- 1944: Oswald Avery (and coworkers) identified DNA as the major carrier of genetic material, responsible for inheritance.
Ribose: (simple) sugar molecule, deoxy-ribose $\rightsquigarrow$ loss of oxygen atom.
Nucleic acid: overall name for DNA and RNA (large biomolecules). Named for their initial discovery in nucleus of cells, and for presence of phosphate groups (related to phosphoric acid).


Ribose


Deoxyribose

## Short historical Introduction

- 1953, Watson \& Crick: 3-dimensional structure of DNA. They inferred the method of DNA replication.
- 2001: first draft of the human genome published by the Human Genome Project and the company Celera.
- Many new developments, such as Next Generation Sequencing, Deep learning etc.



## Base pairs and the DNA



By Madprime (talk contribs) - Own work, CC BY-SA 3.0,

- DNA composed of 4 basic molecules $\rightsquigarrow$ nucleotides.
- Nucleotides are identical up to different nitrogen base: organic molecule with a nitrogen atom that has the chemical properties of a base (due to free electron pair at nitrogen atom).
- Each nucleotide contains phosphate, sugar (of deoxy-ribose type), and one of the 4 bases: Adenine, Guanine, Cytosine, Thymine (A,G,C,T).
- Hydrogen bonds between base pairs $G \equiv C, A=T$.



Sugar

By OpenStax - https://cnx.org/contents/FPtK1zmh@8.25:fEl3C8Ot@10/Preface, CC BY 4.0, https://commons.wikimedia.org/w/index.php?curid=30131206

## The structure of DNA

- DNA molecule is directional due to asymmetrical structure of the sugars which constitute the skeleton: Each sugar is connected to the strand upstream in its 5 th carbon and to the strand downstream in its 3rd carbon.
- DNA strand goes from $5^{\prime}$ to $3^{\prime}$. The directions of the two complementary DNA strands are reversed to one another ( $\rightsquigarrow$ Reversed Complement).


Adapted from https://commons.wikimedia.org/w/index.php?curid=30131206


By Zephyris - Own work, CC BY-SA 3.0, https://commons.wikimedia.org/w/index.php?curid=15027555

## Replication of DNA

Biological process of producing two replicas of DNA from one original DNA molecule. Cells have the distinctive property of division
$\rightsquigarrow$ DNA replication is most essential part for biological inheritance.
Unwinding $\rightsquigarrow$ single bases exposed on each strand.
Pairing requirements are strict $\rightsquigarrow$ single strands are templates for re-forming identical double helix (up to mutations).
DNA polymerase: enzyme that catalyzes the synthesis of new DNA.


## Genes and Chromosomes

- In higher organisms, DNA molecules are packed in a chromosome.
- Genome: total genetic information stored in the chromosomes.
- Every cell contains a complete set of the genome, differences are due to variable expression of genes.
- A gene is a sequence of nu-


By Sponk, Tryphon, Magnus Manske,
https://commons.wikimedia.org/w/index.php?curid=20539140 cleotides that encodes the synthesis of a gene product.

- Gene expression: Process of synthesizing a gene product (often a protein) $\rightsquigarrow$ controls timing, location, and amount.


## The Central Dogma



Transcription: making of an RNA molecule from DNA template. Translation: construction of amino acid sequence from RNA.
$\Rightarrow$ Almost no exceptions ( $\rightsquigarrow$ retroviruses)

## Transcription



1 RNA polymerase (not shown) adds complementary RNA nucleotides to a template DNA strand. The formed RNA strand is identical to the other coding DNA strand, except U is substituted for $T$.


2 Various proteins bil to a sequence AAU near the 3' end of $t$ pre-mRNA molecul 10-30 nucleotides the cleavage and $p$ specificity factor (C) the pre-mRNA.

By Kelvinsong - Own work, CC BY 3.0, https://commons.wikimedia.org/w/index.php?curid=23086203


## Translation

- mRNA molecules are translated by ribosomes: Enzyme that links together amino acids.
- Message is read three bases at a time.
- Initiated by the first AUG codon (codon = nucleotide triplet).
- Covalent bonds (=sharing of electron pairs) are made between adjacent amino acids $\Rightarrow$ growing chain of amino acids
("polypeptide").
- When a "stop" codon (UAA, UGA, UAG) is encountered, translation stops.


RNA
Ribonucleic acid


## Peptide Synthesis

By Boumphreyfr - Own work, CC BY-SA 3.0, https://commons.wikimedia.org/w/index.php?curid=7200200

## The genetic code

Standard genetic code

| 1st base | 2nd base |  |  |  |  |  |  |  | 3rd <br> base |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | U |  | c |  | A |  | G |  |
| U | UUU | (Phe/F) Phenylalanine | $\begin{aligned} & \mathrm{UCU} \\ & \hline \text { UCC } \end{aligned}$ | (Ser/S) Serine | UAU | (Tyr/M) Tyrosine | UGU | (Cys/C) Cysteine | บ |
|  | UUC |  |  |  | UAC |  | UGC |  | c |
|  | UUA | (Leul) Leucine | UCA |  | UAA $A^{[8]}$ | Stop (Ochre) | UGA ${ }^{[8]}$ | Stop (Opal) | A |
|  | UUG |  | UCG |  | UAG ${ }^{[8]}$ | Stop (Amber) | UGG | (TrpW) Tryptophan | G |
| C | CUU |  | ccu | (Pro/P) Proline | CAU | (His/H) Histidine | cGU | (Arg/R) Arginine | U |
|  | cuc |  | CCO |  | CAC |  | CGC |  | c |
|  | CUA |  | CCA |  | CAA | (Gin/Q) Glutamine | CGA |  | A |
|  | CuG |  | CCG |  | CAG |  | CGG |  | G |
| A | AUU | (Iell) Ispleucine | ACU | (Thr/T) Threonine | AAU | (Asn/N) Asparagine | AGU | (Ser/S) Serine | U |
|  | AUC |  | ACC |  | AAC |  | AGC |  | c |
|  | AUA |  | ACA |  | AAA | (Lys/K) Lysine | AGA | (Arg/R) Arginine | A |
|  | Aug ${ }^{[A]}$ | (MetM) Methionine | ACO |  | AAG |  | AGG |  | G |
| G | GUU | (Val/V) Valine | GCu | (Ala/A) Alanine | GAU | (Asp/D) Aspartic acid | GGU | (Gly/G) Glycine | U |
|  | GUC |  | GOC |  | GAC |  | GGC |  | c |
|  | GUA |  | GCA |  | GAA | (Glu/E) Glutamic acid | GGA |  | A |
|  | GUG |  | GOG |  | GAG |  | GGG |  | G |

Wikipedia
Highly redundant: only 20 (or 21) amino acids formed from $4^{3}=64$ possible combinations.







C. Special Cases


Glycine Gly) G

Proline ${ }^{\text {(Prol }} P$


D. Amino Acids with Hydrophobic Side Chain


Phenylalanine (Phe) F Tyrosine (Tyr) Y Tryptophan (Trp) W
 $\mathrm{H}_{2}$
 pK: 10.10
By Dancojocari. https://commons.wikimedia.org/w/index.php?curid=9176441

## Proteins

- Linear polymer of amino acids, linked together by peptide bonds. Average size $\approx 200$ amino acids, can be over 1000 .
- To a large extent, cells are made of proteins.
- Proteins determine shape and structure of a cell. Main instruments of molecular recognition and catalysis.
- Complex structure with four hierarchical levels.

1. Primary structure: amino acid sequence.
2. Different regions form locally regular secondary structures like $\alpha$-helices and $\beta$-sheets.
3. Tertiary structure: packing such structures into one or several 3D domains.
4. Several domains arranged in a quaternary structure.

## Molecular recognition

Interaction between molecules through noncovalent bonding


Crystal structure of a short peptide L-Lys-D-Ala-D-Ala (bacterial cell wall precursor) bound to the antibiotic vancomycin through hydrogen
bonds. By M stone, CC BY-SA 3.0, https://commons.wikimedia.org/w/index.php?curid=2327682

## Catalysis

Increasing the rate of a chemical reaction by adding a substance $\rightsquigarrow$ catalyst.


## Protein Structure: primary to quaternary



Durbin et al., Cambridge University Press
Structure is determined by the primary sequence and their physico-chemical interactions in the medium. Structure determines functionality.

## Mutations

- Mutation: Heritable change in the DNA sequence. Occur due to exposure to ultra violet radiation or other environmental conditions.
- Two levels at which a mutation can take place:
- Point mutation: within a single gene.
- substitution (change of one nucleotide),
- insertion (addition of nucleotides),
- deletion.
- Chromosomal mutation: whole segments interchange, either on the same chromosome, or on different ones.


## Point Mutations

- May arise from spontaneous mutations during DNA replication.
- The rate of mutation increased by mutagens (physical or chemical agent that changes the genetic material).
- Mutagens: Physical (UV-, X-rays or heat), or chemical (molecules misplace base pairs / disrupt helical shape of DNA).



## Importance of Mutations

- Mutations are responsible for inherited disorders \& diseases. Sickle-cell anemia caused by missense point mutation in hemoglobin (in blood cells, responsible for oxygen transport.) Hydrophilic glutamic acid replaced with hydrophobic valine.
$\rightsquigarrow$ deformed red blood cells.
Sequence for Normal Hemoglobin: 6th codon: adenine (A)

| AUG | GUG | CAC | CUG | ACU | CCU | GAG | GAG | AAG | UCU | GCC | GUU | ACU |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| START | Val | His | Leu | Thr | Pro | Glu | Glu | Lys | Ser | Ala | Val | Thr |

Sickle Cell Hemoglobin: $\rightsquigarrow$ thymine (DNA), uracil (RNA)

| AUG | GUG | CAC | CUG | ACU | CCU | GUG | GAG | AAG | UCU | GCC | GUU | ACU |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| START | Val | His | Leu | Thr | Pro | Val | Glu | Lys | Ser | Ala | Val | Thr |

- Mutations are the source of phenotypic variation
$\Rightarrow$ new species and adaption to environmental conditions.


## Sequence Comparison: Motivation

## Basic idea: similar sequences $\rightsquigarrow$ similar proteins.

Protein folding: 30 \% sequence identity $\Rightarrow$ structures similar.


Rout et al., Scientific Reports, vol 8, no 7002 (2018)

## Comparing sequences

Theory: during evolution mutations occurred, creating differences between families of contemporary species.

Missense mutation

Original DNA code for an amino acid sequence.

U.S. National Librany of Medicine

## Comparing sequences

## Comparing two sequences: looking for evidence that they have diverged from a common ancestor by a mutation process.



Thomas Shafee - Own work, CC BY 4.0, https://commons.wikimedia.org/w/index.php?curid=37188728

## Sequence Alignment

Informal definition:
Alignment of sequences $x=x_{1} \ldots x_{n}$ and $y=y_{1} \ldots y_{m}$ :
(i) insert spaces,
(ii) place resulting sequences one above the other so that every character or space has a counterpart.

Example: $\operatorname{ACBCDDDB}$ and CADBDAD. Possible alignments:

$$
\begin{array}{llllllllll}
\mathrm{A} & \mathrm{C} & - & - & \mathrm{B} & \mathrm{C} & \mathrm{D} & \mathrm{D} & \mathrm{D} & \mathrm{~B} \\
& \mid & & & \mid & & \mid & & \mid & \\
- & \mathrm{C} & \mathrm{~A} & \mathrm{D} & \mathrm{~B} & - & \mathrm{D} & \mathrm{~A} & \mathrm{D} & - \\
- & \mathrm{A} & \mathrm{C} & \mathrm{~B} & \mathrm{C} & \mathrm{D} & \mathrm{D} & \mathrm{D} & \mathrm{~B} & \\
& \mid & & \mid & & & \mid & & \\
\mathrm{C} & \mathrm{~A} & \mathrm{D} & \mathrm{~B} & \mathrm{D} & \mathrm{~A} & \mathrm{D} & - & -
\end{array}
$$

## Optimal Alignment

Given: two sequences $x$ and $y$ over alphabet $\mathcal{A}$.
$\mathcal{A}=\{\mathrm{A}, \mathrm{G}, \mathrm{C}, \mathrm{T}\}$ (DNA)
$\mathcal{A}=\{\mathrm{A}, \mathrm{R}, \mathrm{N}, \mathrm{D}, \mathrm{C}, \mathrm{Q}, \mathrm{E}, \mathrm{G}, \mathrm{H}, \mathrm{I}, \mathrm{L}, \mathrm{K}, \mathrm{M}, \mathrm{F}, \mathrm{P}, \mathrm{S}, \mathrm{T}, \mathrm{W}, \mathrm{Y}, \mathrm{V}\}$ (proteins)
Formalizing optimality of an alignment: define

- the costs for substituting a letter by another letter $\Rightarrow$ substitution matrix;
- the costs for insertion $\Rightarrow$ gap penalties.


## The Scoring Model

- Idea: assign a score to each alignment, choose best one.
- Additive scoring scheme: Total score = sum of all scores for pairs of letters + costs for gaps. Implicit assumption:
Mutations at different sites have occurred independently. (In most cases) reasonable for DNA and protein sequences.
- All common algorithms use additive scoring schemes.
- Modeling dependencies is possible, but at the price of significant computational complexities.


## Substitution Matrices

- Expectation: Identities in real alignments are more likely than by chance.
- Derive score for aligned pairs from a probabilistic model.
- Score: relative likelihood that two sequences are evolutionary related as opposed to being unrelated $\rightsquigarrow$ score $=$ ratio of probabilities.
- First assumption: Ungapped alignment, $n=m$.
- $R$ : Random model:

Letter $a$ occurs independently with some frequency $q_{a}$
$\Rightarrow \operatorname{Pr}($ two sequences $)=$ product of probabilities for each letter:

$$
P(x, y \mid R)=\prod_{i} q_{x_{i}} \prod_{i} q_{y_{i}}
$$

## Substitution Matrices

- $M$ (match): aligned pairs occur with joint probability

$$
P(x, y \mid M)=\prod_{i} p_{x_{i} y_{i}}
$$

- Ratio $\rightsquigarrow$ "odds ratio":

$$
\frac{P(x, y \mid M)}{P(x, y \mid R)}=\prod_{i} \frac{p_{x_{i} y_{i}}}{q_{x_{i}} q_{y_{i}}}
$$

- To arrive at an additive scoring system $\rightarrow$ log-odds ratio:

$$
S=\sum_{i} \log \left(\frac{p_{x_{i} y_{i}}}{q_{i} q_{y_{i}}}\right)=\sum_{i} s\left(x_{i}, y_{i}\right)
$$

- $s(a, b)$ : log-likelihood ratio of pair $(a, b)$ occurring as an aligned pair as opposed to an unaligned pair $\rightsquigarrow$ substitution matrix.


## BLOSUM62 substitution matrix

| Ala | 4 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Arg | -1 | 5 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| Asn | -2 | 0 | 6 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| Asp | -2 | -2 | 1 | 6 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| Cys | 0 | -3 | -3 | -3 | 9 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| GIn | -1 | 1 | 0 | 0 | -3 | 5 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| Glu | -1 | 0 | 0 | 2 | -4 | 2 | 5 |  |  |  |  |  |  |  |  |  |  |  |  |  |
| Gly | 0 | -2 | 0 | -1 | -3 | -2 | -2 | 6 |  |  |  |  |  |  |  |  |  |  |  |  |
| His | -2 | 0 | 1 | -1 | -3 | 0 | 0 | -2 | 8 |  |  |  |  |  |  |  |  |  |  |  |
| Ile | -1 | -3 | -3 | -3 | -1 | -3 | -3 | -4 | -3 | 4 |  |  |  |  |  |  |  |  |  |  |
| Leu | -1 | -2 | -3 | -4 | -1 | -2 | -3 | -4 | -3 | 2 | 4 |  |  |  |  |  |  |  |  |  |
| Lys | -1 | 2 | 0 | -1 | -3 | 1 | 1 | -2 | -1 | -3 | -2 | 5 |  |  |  |  |  |  |  |  |
| Met | -1 | -1 | -2 | -3 | -1 | 0 | -2 | -3 | -2 | 1 | 2 | -1 | 5 |  |  |  |  |  |  |  |
| Phe | -2 | -3 | -3 | -3 | -2 | -3 | -3 | -3 | -1 | 0 | 0 | -3 | 0 | 6 |  |  |  |  |  |  |
| Pro | -1 | -2 | -2 | -1 | -3 | -1 | -1 | -2 | -2 | -3 | -3 | -1 | -2 | -4 | 7 |  |  |  |  |  |
| Ser | 1 | -1 | 1 | 0 | -1 | 0 | 0 | 0 | -1 | -2 | -2 | 0 | -1 | -2 | -1 | 4 |  |  |  |  |
| Thr | 0 | -1 | 0 | -1 | -1 | -1 | -1 | -2 | -2 | -1 | -1 | -1 | -1 | -2 | -1 | 1 | 5 |  |  |  |
| Trp | -3 | -3 | -4 | -4 | -2 | -2 | -3 | -2 | -2 | -3 | -2 | -3 | -1 | 1 | -4 | -3 | -2 | 11 |  |  |
| Tyr | -2 | -2 | -2 | -3 | -2 | -1 | -2 | -3 | 2 | -1 | -1 | -2 | -1 | 3 | -3 | -2 | -2 | 2 | 7 |  |
| Val | 0 | -3 | -3 | -3 | -1 | -2 | -2 | -3 | -3 | 3 | 1 | -2 | 1 | -1 | -2 | -2 | 0 | -3 | -1 |  |
|  | Ala | Arg | Asn | Asp | Cys | Gln | Glu | Gly | His | lle | Leu | Lys | Met | Phe | Pro | Ser | Thr | Trp | Tyr |  |

## Gap penalties

Gap penalty types for a gap of length $g$ :

- Linear: $\gamma(g)=-g d$, with $d$ being the gap weight.
- Affine: $\gamma(g)=-d-(g-1) e$, gap-open penalty $d$, gap-extension penalty $e$. Usually $e<d$.
- Convex: e.g. $\gamma(g)=-d \log (g)$. Each additional space contributes less than the previous space.



## Global Alignment: Needleman-Wunsch algorithm

## The Global Alignment problem:

INPUT: two sequences $x=x_{1} \ldots x_{n}$ and $y=y_{1} \ldots y_{m}$.
TASK: Find optimal alignment for linear gap penalties $\gamma(g)=-g d$.
Let $F(i, j)$ be the optimal alignment score of the prefix sequences $x_{1 \ldots i}$ and $y_{1 \ldots j}$. A zero index $i=0$ or $j=0$ refers to an empty sequence. $F(i, j)$ has following properties:

Base conditions: $\quad F(i, 0)=\sum_{k=1}^{i}-d=-i d$

$$
F(0, j)=\sum_{k=1}^{j}-d=-j d, \quad F(0,0)=0 .
$$

Recurrence relation:

$$
\begin{gathered}
\text { for } 1 \leq i \leq n, 1 \leq j \leq m: \\
F(i, j)=\max \left\{\begin{array}{l}
F(i-1, j-1)+s\left(x_{i}, y_{j}\right) \\
F(i-1, j)-d \\
F(i, j-1)-d
\end{array}\right.
\end{gathered}
$$

## Tabular Computation of Optimal Alignment

Starting from $F(0,0)=0$, fill the whole matrix $(F)_{i j}$ :
for $i=0$ or $j=0$, calculate new value from left-hand (upper) value.
for $i, j \geq 1$, calculate the bottom right-hand corner of each square of 4 cells from one of the 3 other cells:

| $\mathrm{F}(0,0)$ <br> $\mathbf{0}$ | $\mathrm{F}(1,0)$ <br> -d | $\mathrm{F}(2,0)$ <br> -2 d |  |
| :---: | :---: | :---: | :---: |
| $\mathrm{F}(0,1)$ <br> $-\mathrm{d} d$ |  |  |  |
| $\mathrm{F}(0,2)$ <br> -2 d |  |  |  |


|  |  |  |  |
| :--- | :--- | :--- | :--- |
|  | $F(i-1, j-1)$ | $F(i, j-1)$ |  |
| $+s\left(\mathbf{x}_{\mathbf{i}}, \mathbf{y}_{\mathbf{j}}\right)$ |  | $-\mathbf{d}$ |  |
|  | $F(\mathrm{i}-1, \mathrm{j})$ | $\mathrm{F}(\mathrm{i}, \mathrm{j})$ |  |
|  | $-\mathbf{d}$ |  |  |
|  |  |  |  |

keep a pointer in each cell back to the cell from which it was derived $\Rightarrow$ traceback pointer.

## Global Alignment: Example

$x=$ HEAGAWGHEE, $y=$ PAWHEAE. Linear gap costs $d=8$. Scoring matrix: BLOSUM50


Durbin et al., Cambridge University Press

## Example: traceback procedure

|  |  | $\begin{aligned} & \mathrm{H} \\ & -8< \end{aligned}$ | $\begin{aligned} & E \\ & --16 \end{aligned}$ | $\begin{aligned} & \text { A } \\ & -24 \end{aligned}$ | $\begin{aligned} & G \\ & -32 \end{aligned}$ | $\begin{aligned} & \text { A } \\ & -40 \end{aligned}$ | $\begin{aligned} & W \\ & -48 \end{aligned}$ | $\begin{aligned} & \mathrm{G} \\ & -56 \end{aligned}$ | $\begin{aligned} & \mathrm{H} \\ & -64 \end{aligned}$ | $\begin{aligned} & \mathrm{E} \\ & -72 \end{aligned}$ | $\begin{aligned} & \mathrm{E} \\ & -80 \end{aligned}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| P | $\begin{array}{r} \uparrow \\ -8 \\ 4 \end{array}$ | $-2$ |  |  | -25 | -33 | -42 | -49 | -57 | -65 | -73 |
| A | -16 | -10 | -3 | -4 | -12 | -20 | -28 | -36 | -44 | -52 | -60 |
| W | -24 | -18 | -11 | -6 | $-7$ | -15 | -5 | -13 | -21 | -29 | -37 |
| H | -32 | -14 | -18 | -13 | -8 | -9 | -13 | -7 | -3 | -11 | -19 |
| E | -40 | -22 | -8 | -16 | -16 | -9 | -12 | -15 | -7 | $\begin{aligned} & 3 \\ & 4 \end{aligned}$ | -5 |
| A | -48 | $-30$ | -16 | -3 | -11 | -11 | -12 | -12 | -15 | $-5$ | 2 |
| E | -56 | $-38$ | -24 | -11 | -6 | -12 | -14 | -15 | -12 | -9 | 1 |
|  | H | E | A | G | A | W | G | H | E | - | E |
|  | - | - | P | - | A | W | - | H | E | A | E |

Add pair of symbols: $\nwarrow:\left(x_{i}, y_{j}\right), \uparrow:\left(-, y_{j}\right), \leftarrow:\left(x_{i},-\right)$

Adapted from Durbin et al., Cambridge University Press

## Time and Space Complexity

Theorem 1. The time complexity of the Needleman-Wunsch algorithm is $O(n m)$. Space complexity is $O(m)$, if only $F(x, y)$ is required, and $O(n m)$ for the reconstruction of the alignment.

## Proof:

Time: when computing $F(i, j)$, only cells $(i-1, j-1),(i, j-1),(i-1, j)$ are examined $\rightsquigarrow$ constant time. There are $(n+1)(m+1)$ cells $\rightsquigarrow O(n m)$ time complexity.

|  |  |
| :---: | :---: |
| $\left.\begin{aligned} & \mathrm{F}(\mathrm{i}-1, \mathrm{j}-1) \\ & +\mathrm{s}\left(\mathbf{x}_{\mathrm{i}}, \mathbf{y}_{\mathrm{j}}\right) \end{aligned} \right\rvert\,$ | $\begin{gathered} \mathrm{F}(\mathrm{i}, \mathrm{j}-1) \\ -\mathbf{d} \end{gathered}$ |
| $\text { F(i-1, } \mathrm{j}_{-\mathrm{d}} \mid$ | $-F(1, j)$ |
|  |  |

Space : row-wise computation of the matrix: for computing row $k$, only row $k-1$ must be stored $\rightsquigarrow O(m)$ space.
Reconstructing the alignment: all traceback pointers must be stored $\rightsquigarrow O(n m)$ space complexity.

## Global Alignment in Linear Space

- Problem: genomic scale sequence analysis: comparing two large genomic sequences: $m, n \approx 10^{6} \Rightarrow$ space complexity $10^{12}$ is clearly unacceptable!
- Solution: linear space algorithms with space complexity $O(m+n)$.
- Basic idea: divide and conquer. Let $u=\left\lfloor\frac{n}{2}\right\rfloor$ be the integer part of $\frac{n}{2}$.
- Let $v$ be a row index such that the cell $(u, v)$ is on the optimal alignment.
- Split dynamic programming problem into two parts:
$(0,0) \rightarrow(u, v)$ and $(u, v) \rightarrow(n, m)$.
Optimal alignment will be concatenation of individual sub-alignments.
- Repeat splitting until until $u=0$ : trivial


Question: how can we find $v$ ?

## Global Alignment in Linear Space

- For $i \geq u$ define $c(i, j)$ such that $(u, c(i, j))$ is on the optimal path from $(1,1) \rightarrow(i, j)$.

- Let $\left(i^{\prime}, j^{\prime}\right)$ be the preceding cell to $(i, j)$ from which $F(i, j)$ is derived. Update $c(i, j)$ as:

$$
c(i, j)= \begin{cases}j & , \text { if } i=u \\ c\left(i^{\prime}, j^{\prime}\right) & , \text { else }\end{cases}
$$



- Local operation $\rightsquigarrow$ need to store only the previous row of $c()$.
- Finally, $v=c(n, m)$.


## Global Alignment in Linear Space：Example

Computing the $c$ matrix for the first step（ $\mathrm{i}=\mathrm{n}=6, \mathrm{j}=\mathrm{m}=4, \mathrm{u}=3$ ）．
The $c$ values are written as subscripts．BLOSUM62，linear gap costs $d=8$ ．

|  |  | $0$ |  | $\begin{aligned} & 1 \\ & \mathrm{H} \end{aligned}$ |  | $\begin{aligned} & 2 \\ & \mathrm{E} \end{aligned}$ |  | $\begin{aligned} & 3 \\ & \mathrm{~A} \end{aligned}$ |  | $\begin{aligned} & 4 \\ & \mathrm{G} \end{aligned}$ |  | $\begin{aligned} & 5 \\ & \mathrm{~A} \end{aligned}$ |  | $\begin{gathered} 6 \\ \text { W } \end{gathered}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | － | 0 | $\leftarrow$ | －8 | $\leftarrow$ | －16 | $\leftarrow$ | $-24_{0}$ | $\leftarrow$ | $-32_{0}$ | $\leftarrow$ | $-40_{0}$ | $\leftarrow$ | $-48_{0}$ |
|  |  | $\uparrow$ | $\nwarrow$ |  | 「 |  | 「 |  |  |  | $\nwarrow$ |  |  |  |
| 1 | P | －8 |  | －2 |  | －9 |  | $-17_{1}$ | $\leftarrow$ | $-25_{1}$ |  | $-33_{0}$ | $\leftarrow$ | $-41{ }_{0}$ |
|  |  | $\uparrow$ | 「 | $\uparrow$ | 「 |  | 「 |  |  |  | $\nwarrow$ |  |  |  |
| 2 | A | －16 |  | －10 |  | －3 |  | $-42$ | $\leftarrow$ | $-12_{2}$ |  | $-20_{1}$ | $\leftarrow$ | $-28_{1}$ |
|  |  | $\uparrow$ |  | $\uparrow$ |  |  | 「 |  | 「 |  | $\nwarrow$ |  | 「 |  |
| 3 | W | －24 |  | －18 |  | －11 |  | $-63$ |  | $-72$ |  | $-15_{2}$ |  | $-5_{1}$ |
|  |  | $\uparrow$ | $\nwarrow$ |  | 「 |  | 「 |  | 「 |  | $\nwarrow$ |  |  | $\uparrow$ |
| 4 | H | －32 |  | －14 |  | －18 |  | $-13_{4}$ |  | $-83$ |  | $-9_{2}$ |  | $-13_{1}$ |

Every $c(i, j)$ defines a row index $v$ such that $(u, c(i, j))$ is on the optimal path from $(1,1)$ to $(i, j) \rightsquigarrow v=c(6,4)=1$ ，so $(3,1)$ is our desired element on the optimal path form $(1,1)$ to $(6,4)$ ．

## Local Alignments

> The Local Alignment problem:
> INPUT: two sequences $x=x_{1}, \ldots, x_{n}$ and $y=y_{1}, \ldots, y_{m}$.
> TASK: find subsequences $a$ of $x$ and $b$ of $y$,
> whose similarity (=optimal global alignment score) is maximal
> (over all such pairs of subsequences).
> Assume linear gap penalties $\gamma(g)=-g d$.

Subsequence = contiguous segment of a sequence.
Consider first a simpler problem by fixing the endpoint of the subsequences at index pair $(i, j)$ :
Local suffix alignment problem: given $x, y, i, j$, find suffixes $\alpha$ of $x_{1, \ldots, i}$ and $\beta$ of $y=y_{1, \ldots, j}$ such that their global alignment score is maximal.

$$
(x_{1}, \ldots, \underbrace{x_{k}, \ldots, x_{i}}_{\alpha}), \quad(y_{1}, \ldots, \underbrace{y_{l}, \ldots, y_{j}}_{\beta})
$$

## Local suffix alignments

Consider global alignment path to cell $(i, j)$. Where to start? Intuition: Indices $(k, l)$ found by following the path back to $(0,0)$, but stopping at the first negative value.


Remark: If we consider all solutions (i.e. for all $(i, j)$ pairs), we look at all possible subsequences (no restrictions on $\alpha, \beta$ )

Maximal solution of local suffix alignment over all pairs $(i, j)$ = solution of local alignment problem.

## Smith-Waterman Algorithm

$F(i, j)$ : optimal local suffix alignment for indices $i, j$.
Global alignment with one modification:
Prefixes whose scores are $\leq 0$ are discarded
$\rightsquigarrow$ alignment can start anywhere.

Recurrence relation:

$$
F(i, j)=\max \left\{\begin{array}{l}
0 \\
F(i-1, j-1)+s\left(x_{i}, y_{j}\right) \\
F(i-1, j)-d \\
F(i, j-1)-d
\end{array}\right.
$$

Finally, find indices $i^{*}$ and $j^{*}$ after which the similarity only decreases. Stop the alignment there.

$$
F\left(i^{*}, j^{*}\right)=\max _{i, j} F(i, j)
$$

## Traceback...

...starts at highest value until a cell with 0 is reached.


Adapted from Durbin et al., Cambridge University Press

## Local vs. Global Alignment: Biological Considerations

- Many proteins have multiple domains, or modules.
- Some domains are present (with high similarity) in many other proteins
- Local alignment can detect similar regions in otherwise dissimilar proteins.


Durbin et al., Cambridge University Press

## Other gap models

- So far: linear gap model. Not ideal for biological sequences, since it penalizes additional gap steps as much as the first. But in reality: When gaps do occur, they are often longer than one character.

```
HBA_HUMAN GSAQVKGHGKKVADALTNAVAHV---D--DMPNALSALSDLHAHKL
    ++ ++++H+ KV + +A ++ +L+ L+++H+ K
```

LGB2_LUPLU NNPELQAHAGKVFKLVYEAAIQLQVTGVVVTDATLKNLGSVHVSKG

Durbin et al., Cambridge University Press

- For a general gap cost function $\gamma(g)$, we can still use the standard dynamic programming recursion with slight modifications:

$$
F(i, j)=\max \left\{\begin{array}{l}
F(i-1, j-1)+s\left(x_{i}, y_{j}\right) \\
F(k, j)+\gamma(i-k), \quad k=0, \ldots, i-1 \\
F(i, k)+\gamma(j-k), \quad k=0, \ldots, j-1
\end{array}\right.
$$

- Problem: requires $O\left(n^{3}\right)$ operations to align two sequences of length $n$, rather than $O\left(n^{2}\right)$. Why? $\rightsquigarrow$ exercises...


## Alignment with affine gap costs

For affine gap costs, $\gamma(g)=-d-(g-1) e$, there exists a solution: Modify recurrence by introducing another two "states". Denote by

- $M(i, j)$ the best score given that $x_{i}$ is aligned to $y_{j}$,
- $I_{x}(i, j)$ the best score given that $x_{i}$ is aligned to a gap,
- $I_{y}(i, j)$ the best score given that $y_{j}$ is aligned to a gap.

$$
\begin{aligned}
& M(i, j)=\max \left\{\begin{array}{l}
M(i-1, j-1)+s\left(x_{i}, y_{j}\right) \\
I_{x}(i-1, j-1)+s\left(x_{i}, y_{j}\right) \\
I_{y}(i-1, j-1)+s\left(x_{i}, y_{j}\right)
\end{array}\right. \\
& I_{x}(i, j)=\max \left\{\begin{array}{l}
M(i-1, j)-d \\
I_{x}(i-1, j)-e
\end{array}\right. \\
& I_{y}(i, j)=\max \left\{\begin{array}{l}
M(i, j-1)-d \\
I_{y}(i, j-1)-e
\end{array}\right.
\end{aligned}
$$



## Example FSA alignment



FSA alignment corresponds to path through states.
Probabilistic version $\rightsquigarrow$ Hidden Markov models (next chapter)

## Scoring Matrices Revisited: the PAM family

- PAM = Point Accepted Mutations.
(Dayhoff 1978, Atlas of Protein Sequence and Structure, Vol 5.)
- Accepted means that a mutation did not change the function of a protein, or the change was beneficial to the organism.
- PAM matrices are based on global alignments of closely related proteins.
- PAM-1 is the matrix calculated from comparisons of sequences (trusted data!) with no more than 1\% divergence (one mutation per 100 amino acids).
- Other PAM matrices are extrapolated from PAM-1.


## Constructing PAM

## Protein sequences in 71 families, at least $85 \%$ identical. Multiple alignment:

```
KAPPA
    1 HUMAN EU
    2 ~ M O U S E ~ M O P C ~ 2 1 ~
    3 QAT S211
    4 84 RA881T 4135
    5 B9 RA881T
LAMBDA
    6 ~ H U M A N ~ S H
    7 PIG
    8 I MOUSE MOPC 104E
    9 2 MOUSE MOPC 315
```

```
/T - V A A P S V F I F P P S D E Q - L K S - G T A S V V C L L N N F Y P - R E - A
/A - DAAP TVSI F P P S S E Q - L T S - GGASVV CF LNN N Y P - K D - I
/A - NAAP TVSIFPPST Z Z - LA T - GGASVVC LMN K.FYP - R.D - I
/D - PVAP TV L I F P P A A D Q - VA T - G TV T IV V VAN K Y F P - - D - V
/D P P I A P T V L L F P P S A D Q - L T T - Z T V T I V C V A N K F R P - D D - I
/Q P KAAP S V T L F P P S S E E - L Q A - N K A T L V C L I S D F Y P - G A - V
/Q P K A A P T V N L F P P S S E E - L G T - N K A T L V C L I S D F Y P - G A - V
/Q P K S S P S V T L F P P S S E E - L T E - N K A T L V C T I T O F Y P - G V - V
/Q P K S T P T L T V F P P S S E E - L K E - N K.A T L V C.LI S N F S P - G S - (V
```


## Constructing PAM

Build Phylogenetic Tree:


A conceptual phylogenetic tree. Leaves: Four observed proteins. Inner nodes: Inferred ancestors.

Matrix of Replacements

|  | A | B | C | D | G | H | I | J |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| A |  |  | 1 | 1 |  |  |  |  |
| B |  |  | 1 | 1 |  |  |  |  |
| C | 1 | 1 |  |  |  |  |  |  |
| D | 1 | 1 |  |  |  |  |  |  |
| G |  |  |  |  |  |  | 1 |  |
| H |  |  |  |  |  |  |  | 1 |
| I |  |  |  |  | 1 |  |  |  |
| J |  |  |  |  |  | 1 |  |  |

Matrix of accepted point mutations derived from the tree.

## Constructing PAM

Cumulative data from (Dayhoff, M.O., Schwartz, R. and Orcutt, B.C. (1978). A model of Evolutionary Change in Proteins. Atlas of protein sequence and structure (volume 5, supplement 3 ed.) pp. 345358)

|  | ala | arg | asn | asp | cys | gln | glu | gly | his | ile | leu | lys | met |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A |  |  |  |  |  |  |  |  |  |  |  |  |  |
| R | 30 |  |  |  |  |  |  |  |  |  |  |  |  |
| N | 109 | 17 |  |  |  |  |  |  |  |  |  |  |  |
| D | 154 | 0 | 532 |  |  |  |  |  |  |  |  |  |  |
| C | 33 | 10 | 0 | 0 |  |  |  |  |  |  |  |  |  |
| Q | 93 | 120 | 50 | 76 | 0 |  |  |  |  |  |  |  |  |
| E | 266 | 0 | 94 | 831 | 0 | 422 |  |  |  |  |  |  |  |
| G | 579 | 10 | 156 | 162 | 10 | 30 | 112 |  |  |  |  |  |  |
| H | 21 | 103 | 226 | 43 | 10 | 243 | 23 | 10 |  |  |  |  |  |
| I | 66 | 30 | 36 | 13 | 17 | 8 | 35 | 0 | 3 |  |  |  |  |
| L | 95 | 17 | 37 | 0 | 0 | 75 | 15 | 17 | 40 | 253 |  |  |  |
| K | 57 | 477 | 322 | 85 | 0 | 147 | 104 | 60 | 23 | 43 | 39 |  |  |
| M | 29 | 17 | 0 | 0 | 0 | 20 | 7 | 7 | 0 | 57 | 207 | 90 |  |
| F | 20 | 7 | 7 | 0 | 0 | 0 | 0 | 17 | 20 | 90 | 167 | 0 | 17 |
| P | 345 | 67 | 27 | 10 | 10 | 93 | 40 | 49 | 50 | 7 | 43 | 43 | 4 |
| S | 772 | 137 | 432 | 98 | 117 | 47 | 86 | 450 | 26 | 20 | 32 | 168 | 20 |
| T | 590 | 20 | 169 | 57 | 10 | 37 | 31 | 50 | 14 | 129 | 52 | 200 | 28 |
| W | 0 | 27 | 3 | 0 | 0 | 0 | 0 | 0 | 3 | 0 | 13 | 0 | 0 |
| Y | 20 | 3 | 36 | 0 | 30 | 0 | 10 | 0 | 40 | 13 | 23 | 10 | 0 |
| V | 365 | 20 | 13 | 17 | 33 | 27 | 37 | 97 | 30 | 661 | 303 | 17 | 77 |

Numbers of accepted point mutations (x10) accumulated from closely related sequences.

## Constructing PAM: formal derivation

- $f_{A B}$ : frequency of $A$ (in ancestor) replaced by $B$ (in descendant). Assumption: $f_{A B}=f_{B A}$
- $f_{A}=\sum_{B \neq A} f_{A B}$ : number of observations that $A$ is involved in.
- $f=\sum_{A} f_{A}$ : total number of mutations observed.
- $P(B \mid A, t)$ : probability that $A$ is substituted by $B$ in time $t$. One time unit $=$ one "generation" $\Rightarrow P(B \mid A, t=1)=f_{A B} / f_{A}$
- $m_{A}$ : relative mutability of $A=$ likelihood that $A$ is involved in a mutation
$=\frac{\#(\text { mutations } A \text { is involved in })}{\text { total number of mutations } \cdot \text { prob. that a given character is } A}$

$$
\Rightarrow \quad m_{A}=\frac{f_{A}}{f \cdot P_{A}} .
$$

## Constructing PAM: formal derivation (cont'd)

- $M_{A B}$ : probability that $A$ mutates to $B$ in $t=1$ : $P(B \mid A, t=1$, match $)$ Product of mutability of $A$ and probability that given $A$ has mutated, it has mutated to $B$ in time $t=1$.

$$
M_{A B}=P(B \mid A, t=1) m_{A}=\frac{f_{A B}}{f_{A}} m_{A}=\frac{f_{A B}}{f \cdot P_{A}} .
$$

- Expected number of mutations in one time unit:

$$
\sum_{A} P_{A} \sum_{B \neq A} M_{A B}=\sum_{A} P_{A} \sum_{B \neq A} \frac{f_{A B}}{f P_{A}}=\sum_{A} \frac{f_{A}}{f}=1 .
$$

- We want to set $t=1$ when the expected number of mutations is $1 \%$ : $\rightsquigarrow$ we rescale $M_{A B} \leftarrow 0.01 \cdot M_{A B}$.

Model assumption: constant evolutionary clock!

## Constructing PAM: formal derivation (cont'd)

- How to compute the diagonal elements?

Probability that $A$ does not mutate:

$$
\begin{aligned}
& \sum_{B \neq A} M_{A B}+M_{A A} \stackrel{!}{=} 1 \\
& \quad \Rightarrow \quad M_{A A}=1-\sum_{B \neq A} M_{A B}=1-0.01 \cdot \frac{f_{A}}{f P_{A}}=1-0.01 \cdot m_{A} .
\end{aligned}
$$

- $M$ is the PAM-1 matrix, i.e. the mutation probability matrix for $t=1$.
- The log-odd scores corresponding to PAM-1 are

$$
s_{A B}=\log \frac{P_{A} \overbrace{M_{A B}}^{P(B \mid A, t=1, \text { match })}}{P_{A} P_{B}}=\log \frac{P(A, B \mid \text { match })}{P(A, B \mid \text { random })} .
$$

## Constructing PAM: formal derivation (cont'd)

- To obtain transition matrices for $t=n$, we multiply $M(t=1)$ by itself $n$ times:

$$
M(t=n)=M^{n}(t=1) .
$$

- $M(t=2)_{A B}$ is the probability that $A$ is substituted by $B$ through an intermediate character $C$.
- Values of $t=40,120,250$ are commonly used.


## The BLOSUM family

- BLOSUM matrices are based on local alignments from protein families in the BLOCKS database.
- Original paper: (Henikoff S \& Henikoff JG, 1992; PNAS 89:10915-10919).
- BLOSUM 62 is a matrix calculated from comparisons of sequences with at least 62\% similarity.
- All BLOSUM matrices are based on observed alignments. They are not extrapolated from comparisons of closely related proteins.

Relationship between BLOSUM and PAM:

| BLOSUM 80 | BLOSUM 62 | BLOSUM 45 |
| :--- | :--- | :--- |
| PAM 1 | PAM 120 | PAM 250 |
|  |  |  |
| Less divergent |  | More divergent |

